

## 4.1 Fibonacci

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12:34 PM

Famous sequence of numbers

0 1 1 2 3 5 8 13 21 34 55 .....

Formula

$$F(n) = F(n-1) + F(n-2) \quad \text{-- recurrence.}$$

$$F(0) = 0$$

$$F(1) = 1$$

} initial conditions

Pseudocode

FUNCTION fibo(n)

IF  $n \leq 1$  THEN

RETURN n

ELSE RETURN fibo(n-1) + fibo(n-2)

Analyse:

basic\_operation. +

Count the number of additions.

$$A(n) = A(n-1) + 1 + A(n-2) \quad \text{-- recurrence}$$

$$A(0) = 0$$

$$A(1) = 0$$

} initial conditions

Note: not easy to solve using backward substitution

So we need to introduce more tools.



### • TOOLS : Linear Second-Order Recurrences

- DEF: linear Second-order Recurrences are Recurrences of the form:

$$a \cdot f(n) + b \cdot f(n-1) + c \cdot f(n-2) = g(n)$$

where: a b c are real numbers

$$a \neq 0$$

"second order"  $f(n-2)$  is "two steps" behind  $f(n)$

two types  $\begin{cases} \text{homogeneous} & \text{if } g(n) = 0 \text{ for all } n \\ \text{inhomogeneous} \end{cases}$

- DEF: the characteristic equation of a homogeneous linear second-order recurrence is:

$$a \cdot r^2 + b \cdot r + c = 0 \quad (r \text{ is the variable})$$

E.g.  $a=13$   
 $b=1$   
 $c=3$   $\cdot 13 \cdot f(n) + 1 \cdot f(n-1) + 3 \cdot f(n-2) = 0$

the characteristic equation is:  $13 \cdot r^2 + 1 \cdot r + 3 = 0$

- THEOREM: from the roots of a characteristic equation we can solve a homogeneous linear second-order recurrence.

Let  $r_1$   $r_2$  be the roots of a characteristic equation:

Case 1:  $r_1$   $r_2$  are real and distinct.  
Then the solution for the recurrence is

$$f(n) = \alpha \cdot r_1^n + \beta \cdot r_2^n$$

where  $\alpha$   $\beta$  are some real numbers.

Case 2:  $r_1$   $r_2$  are equal  
then the solution for the recurrence is

$$f(n) = \alpha \cdot r^n + \beta \cdot n \cdot r^n$$

where  $\alpha$   $\beta$  are some real numbers.

Case 3:  $r_1$   $r_2$  are distinct complex numbers  
then the solution for the recurrence is

$$r_{1,2} = u \pm i \cdot v$$

$$f(n) = \gamma^n \cdot [\alpha \cdot \cos(n \cdot \theta) + \beta \cdot \sin(n \cdot \theta)]$$

where:

$$\gamma = \sqrt{u^2 + v^2}$$

$$\theta = \arctan\left(\frac{v}{u}\right)$$

$\alpha$   $\beta$  some real numbers

Where:

$$\gamma = \sqrt{u^2 + v^2}$$

$$\theta = \arctan\left(\frac{v}{u}\right) \quad \alpha \quad \beta \quad \text{Some real numbers.}$$

## • REVIEW: Quadratic Formula

$$a \cdot r^2 + b \cdot r + c$$

$$r = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

## • Back to Fibonacci:

Formula

$$F(n) = F(n-1) + F(n-2) \quad \text{— recurrence.}$$

$$F(0) = 0$$

$$F(1) = 1$$

} initial conditions

$$F(n) - F(n-1) - F(n-2) = 0$$

$$a = 1$$

$$b = -1$$

$$c = -1$$

homogeneous  
linear  
second-order  
recurrence.

characteristic equation:

$$1 \cdot r^2 - 1 \cdot r - 1 = 0$$

with roots:

$$r_{1,2} = \frac{1 \pm \sqrt{1 - 4(1 \cdot -1)}}{2(1)} = \frac{1 \pm \sqrt{5}}{2}$$

then: from Case 1:

$$F(n) = \alpha \cdot \left(\frac{1+\sqrt{5}}{2}\right)^n + \beta \cdot \left(\frac{1-\sqrt{5}}{2}\right)^n$$

Now: Lets find  $\alpha$   $\beta$ :

$$F(0) = 0 = \alpha \cdot \left(\frac{1+\sqrt{5}}{2}\right)^0 + \beta \cdot \left(\frac{1-\sqrt{5}}{2}\right)^0$$

$$0 = \alpha + \beta$$

$$\alpha = -\beta$$

$$\beta = -\alpha$$

$$F(1) = 1 = \alpha \cdot \left(\frac{1+\sqrt{5}}{2}\right)^1 + \beta \cdot \left(\frac{1-\sqrt{5}}{2}\right)^1$$

- $F(1) = 1 = \alpha \cdot \left(\frac{1+\sqrt{5}}{2}\right)' + \beta \cdot \left(\frac{1-\sqrt{5}}{2}\right)'$

Substitute  $\beta$

$$F(1) = 1 = \alpha \cdot \left(\frac{1+\sqrt{5}}{2}\right)' - \alpha \left(\frac{1-\sqrt{5}}{2}\right)'$$

$$1 = \alpha \left[ \frac{1+\sqrt{5}}{2} - \frac{1-\sqrt{5}}{2} \right]$$

$$= \alpha \left[ \frac{1+\sqrt{5} - 1 + \sqrt{5}}{2} \right]$$

$$= \alpha \left[ \frac{2 \cdot \sqrt{5}}{2} \right]$$

$$1 = \alpha [\sqrt{5}]$$

$$\alpha = \frac{1}{\sqrt{5}} \quad \beta = -\frac{1}{\sqrt{5}}$$

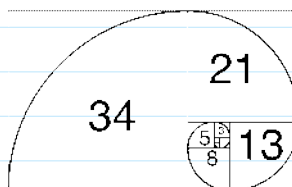
SO: closed formula

$$F(n) = \frac{1}{\sqrt{5}} \cdot \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \cdot \left(\frac{1-\sqrt{5}}{2}\right)^n$$

note:

$$\frac{1+\sqrt{5}}{2} = 1.61803\ldots = \phi \quad \text{"the golden ratio"}$$

$$\text{Let } \hat{\phi} = \frac{-1}{\phi}$$



Then:

$$F(n) = \frac{1}{\sqrt{5}} (\phi^n - \hat{\phi}^n) \in O(\phi^n)$$

Side Note:

You can compute  $F(n) \approx \frac{1}{\sqrt{5}} \cdot \phi^n$  rounded to

Side Note:

you can compute  $F(n) \cong \frac{1}{\sqrt{5}} \cdot \phi^n$  rounded to nearest integer.

• Back to the recursive algorithm fibo()

- $A(n) = A(n-1) + 1 + A(n-2)$  — recurrence
  - $A(0) = 0$
  - $A(1) = 0$
- } initial conditions

- $A(n) - A(n-1) - A(n-2) = \underline{1}$  . it's inhomogeneous

Trick:

Rewrite formula to make it homogeneous.

$$\bullet [A(n)+1] - [A(n-1)+1] - [A(n-2)+1] = 0$$

$$\begin{aligned} A(n)+\cancel{1} - A(n-1) - \cancel{1} - A(n-2) - 1 &= 0 \\ A(n) - A(n-1) - A(n-2) &= 1 \end{aligned}$$

$$\text{Let } B(n) = A(n) + 1$$

$$\bullet B(n) - B(n-1) - B(n-2) = 0$$

$$\begin{aligned} \text{coefficients } a &= 1 \\ b &= -1 \\ c &= -1 \end{aligned}$$

$$\begin{aligned} \text{characteristic equation: } ar^2 + br + c \\ r^2 - r - 1 = 0 \end{aligned}$$

But, Let's take a shortcut

$$B(0) = A(0) + 1 = 1$$

$$B(1) = A(1) + 1 = 1$$

$$B(2) = 2$$

$$B(3) = 3$$

$$B(4) = 5$$

$$B(5) = 8$$

$$B(6) = 13$$

$$B(7) = 21$$

$$B(n) = B(n-1) + B(n-2)$$

$$B(n) = F(n+1)$$

$$\begin{aligned} B(6) &= 13 \\ B(7) &= 21 \\ &\vdots \end{aligned}$$

$$B(n) = F(n+1)$$

So:

$$B(n) = \frac{1}{\sqrt{5}} (\phi^{n+1} - \hat{\phi}^{n+1})$$

$$A(n)-1 = \frac{1}{\sqrt{5}} (\phi^{n+1} - \hat{\phi}^{n+1})$$

$$A(n) = \frac{1}{\sqrt{5}} (\phi^{n+1} - \hat{\phi}^{n+1}) - 1$$

Then

$$A(n) \in \Theta(\phi^n)$$

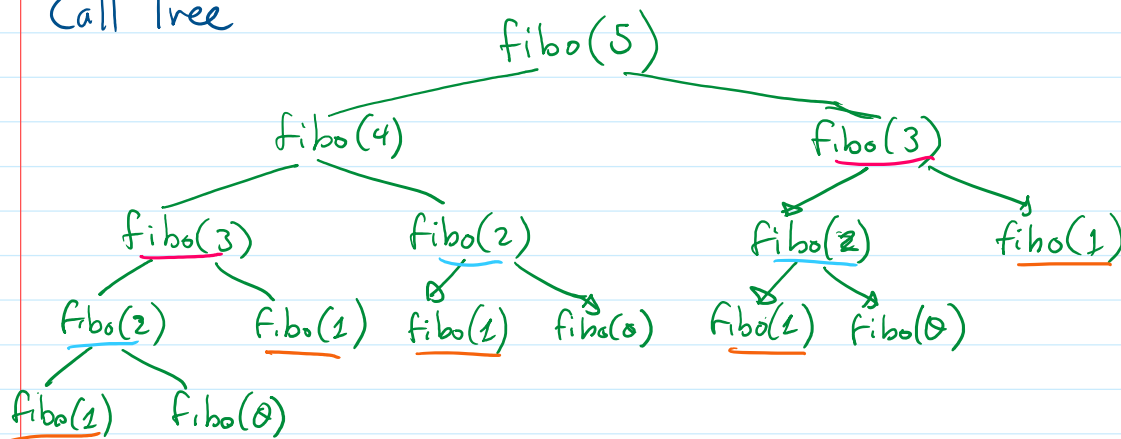
- ANOTHER VIEW : The call tree of fibo()

Pseudocode

```

FUNCTION fibo(n)
  IF n ≤ 1 THEN
    RETURN n
  ELSE RETURN fibo(n-1) + fibo(n-2)
  
```

Call Tree



recursion is powerfull, but its succinctness may hide inefficiencies.

Lets rewrite fibo()

FUNCTION FibArray( F[0...n])

F[0] ← 0

F[1] ← 1

FOR i ← 2 TO n DO

F[i] ← F[i-1] + F[i-2]

RETURN F

0	1	2	3	4	5	6	7	8	9	
0	1	1	2	3	5	8	13	21	34	...

Analysis

basic operation +

$$A(n) = \sum_{i=2}^n 1 = n-2 \in \Theta(n) \text{ linear.}$$

EXTRA:

There is a  $\Theta(\log n)$  way to compute fibonacci, based on:

$$\begin{bmatrix} F(n-1) & F(n) \\ F(n) & F(n+1) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^n$$

and an efficient way to compute powers of matrices

—o— EOF.