

Mathematical Languages.

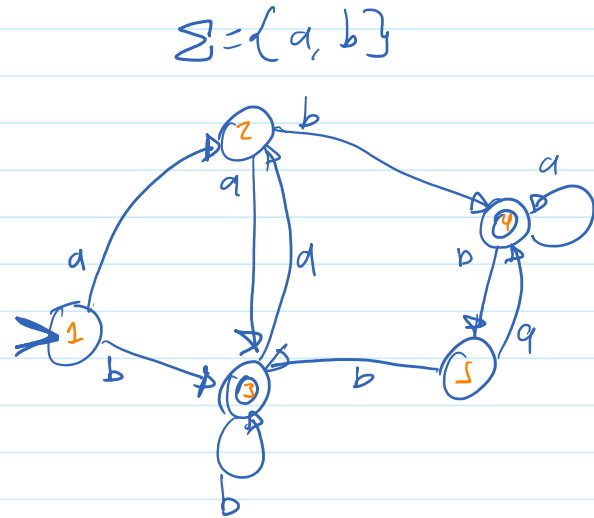
Problems $\left\{ \begin{array}{l} \text{Specification - regex} \\ \text{Recognition. } \underline{i}w \in L? \end{array} \right.$

()
[]
{ }
: ?
! !

• Automata.

- Set of states
 - ↳ initial state
 - ↳ "accept" states

- Transition function
function from state, symbol \rightarrow state.
symbols from Σ



How to recognize a Language?

- Given a word w we start at the initial state and follow the transition function for each symbol from w in order.

$w = aabbaaba$
 $\wedge \wedge \wedge \wedge \wedge \wedge \wedge \wedge$
 $\rightarrow 4$

- if you end in an accept state

$w \in L!$

otherwise

$w \notin L!$

$L = \begin{cases} ab & bab \\ b & bbbbbb \end{cases}$

~~THEOREM~~ For every Language L specified by a reg-ex. There exists an automata that recognizes L

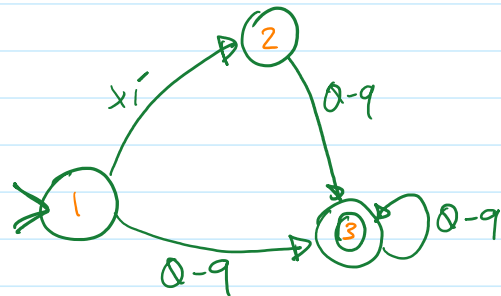
E.G.

Integer Constants

$(+|-)?[0-9]^+$

$\Sigma = \text{ASCII}$

0 -9 27
+82

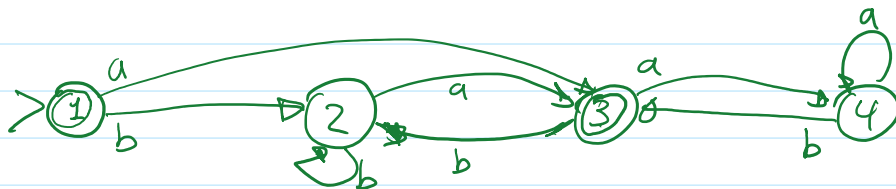


Note: There exist an algorithm that takes a regex and constructs an automata.

• Encoding an automata.

- Loop over a switch statement.
- every case of switch statement will correspond to a state in the automata.
- the body of each switch statement should correspond to the relevant transition function.

E.G.



$L = \{ a, ba, bba \}$

Pseudocode

```

FUNCTION recognize( string s ) : BOOLEAN
VAR state, I : INTEGER
    c : CHAR
BEGIN
  
```

```

state := 1; i := 0

WHILE i < length(s) DO
  c := s[i]
  CASE state OF
    1 : IF c = 'a' THEN state := 3
        ELSE state := 2
    2 : IF c = 'a' THEN state := 3
        ELSE state := 2
    3 : IF c = 'a' THEN state := 4
        ELSE state := 2
    4 : IF c = 'a' THEN state := 4
        ELSE state := 3
  END
  i := i + 1
END
IF state = 3 OR state = 1 THEN
  RETURN True
ELSE RETURN False
END

```

We could:

- add fail / sink states
- return early
- turn switch into lookup-table.

this algorithm is the basis of our
 "Lexical analyzer"



C++: an octopus made by nailing extra legs onto a dog. – Steve Taylor

— EOF —