

"Static" :- as in the text
"Dynamic" :- as the program runs.

TYPE BINDING

When does a variable get its type?

Static type Binding

Type is bound somewhere in the text.

- Explicit statement

C++ int x; string s;

Pascal
VAR
X: INTEGER;
S: STRING;

- Implicit Declaration

FORTRAN: I J K are integers
everything else is a float.
(unless otherwise specified)

Perl:
name carries type.

@x array
%x hashtable / Dictionary
\$x scalar.

- Static Type Deduction

Type is deduced by initialization.

Go

VAR float x
:
x = 3.14

VAR x := 3.14

↑ apply type deduction.

Dynamic type Binding

- Type is bound when a value is assigned, as the program runs.

- type is bound when a value is assigned, as the program runs.
- type can change, by another assignment.

Python JavaScript Ruby Lua

■ SCOPE

- The scope of a variable is the range of statements over which the variable is visible

- Global:

BASIC

- Static Scope.

range is determined by program text.

C/C++

scope is based on "blocks"
 { block }

e.g.

```
int z; //global
```

```
{ int foo()
{
```

```
  int z;
```

```
  ...
```

```
  }
```

```
  int z;
```

```
  ...
```

```
  }
```

```
}
```

← local variables mask global variables

← inner scope masks outer scope.

Pascal

Scope is Module or Function based.

```
VAR z: INTEGER
```

```
FUNCTION foo():INTEGER:
```

```
  VAR z: INTEGER
```

```
  BEGIN
```

```
    ...
```

```
  END.
```

← Declarations allowed here only.

the scope of z is the whole function.

- How Does the symbol table handle nested scopes?

- How Does the symbol table handles nested scopes?
 - Split symbol table into "frames"
 - new scope:- push a new frame
 - end scope:- pop a frame
 - Symbol table is now a stack

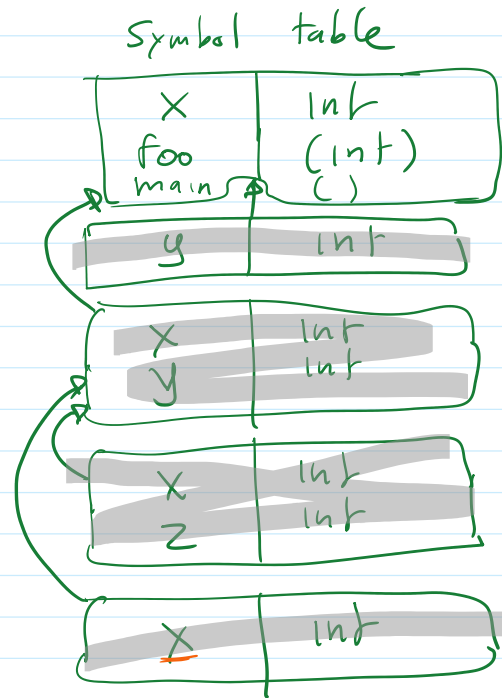
E.G C++

```

→ int x;
→ void foo ( int y )
→ {
→   x = y * 10;
→ }

→ int main()
→ {
→   int x, y;
→   cin >> x >> y;
→   if ( x < y ) {
→     int x, z;
→     cin >> z;
→     x = y + z;
→   }
→   for( int x=0; x<y; x++) {
→     cout << x,
→   }
→   cout << x;
→ }

```



• Dynamic Scope.

- Used in some esoteric programming languages.
- '70:- important among the LISP community
- When is a variable visible?:- Depends on program execution
 - If you have a function call: fun1() calls fun2() then the variables from fun1() are available in fun2().

E.G.

Imagine C++ w/ Dynamic Scope

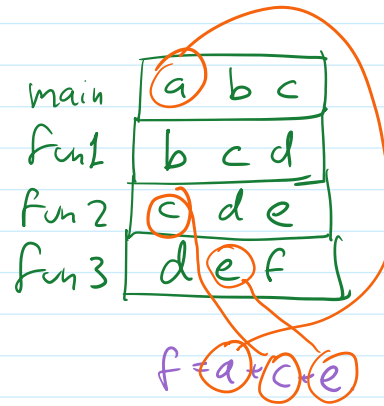
```
void main ()
{
  int a, b, c
  ... ..
}
```

```
void fun1 ()
{
  int b, c, d
  ... ..
  b = a + 1
}
```

```
void fun2 ()
{
  int c, d, e
  ... ..
}
```

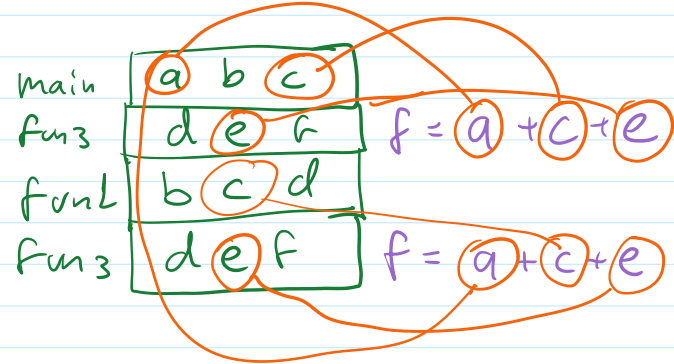
```
void fun3 ()
{
  int d, e, f
  ... ..
  f = a + c + e
}
```

Suppose:
main
 fun1
 fun2
 fun3



Suppose:

main
 fun3
 fun1
 fun3



E.5 #2

"C++" w/ Dynamic scope.

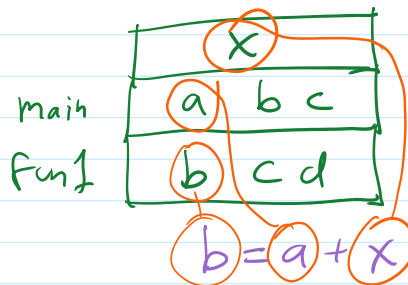
```
int x

void main ()
{
  int a, b, c
  ... ..
}
```

```
void fun1 ()
{
  int b, c, d
  ... ..
  b = a + x
}
```

Suppose
main
 fun1

call stack/symbol table.



}

—o—o— Eof.