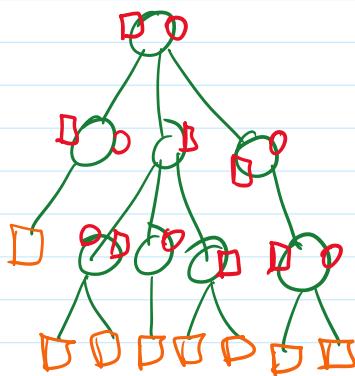


- Main idea:

What if we could attach semantic information to our parse tree.

extra information on the intermediate nodes.



- ATTRIBUTE GRAMMAR:

- expand parse tree with semantic info:
  - D. Knuth and Wagner.

Idea:

- attach to each node a collection of attributes.

- for each grammar symbol  $X$   
a set of attributes  $\text{Att}(X)$

- for each grammar rule  $R$   
a collection of rules: attribute rules  
that assign values to attributes of symbols in  $R$ .

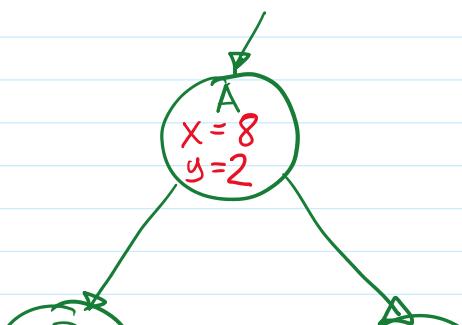
E.G. #1

attributes  $x \ y$

Grammar rule:  $A \rightarrow BC$  ... Attribute rules:

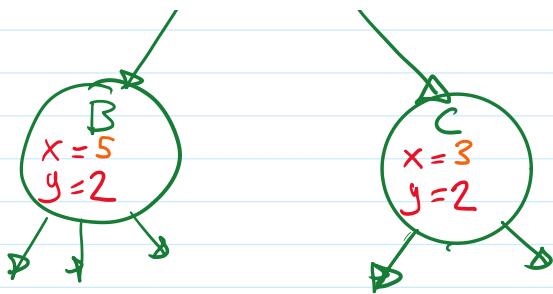
$$A \rightarrow BC$$

$$\begin{aligned} A.x &\leftarrow B.x + C.x \\ A.y &\leftarrow 2 \\ B.y &\leftarrow A.y \\ C.y &\leftarrow B.y \end{aligned}$$



Synthesized:  
 $A.x \ A.y$

Inherited:



Inherited:  
B.y  
C.y

- Types of attributes:

- "Synthesized" Attributes:

- value depends on the attribute values of a node's children.
- info flows from bottom to top of the tree.

- "Inherited" Attributes

- value depends on the values of attributes of node's siblings or parents
- info flows {from top to bottom} of the tree  
[- Sideways]

E.G. #2

attributes: type exptype.

$A \rightarrow \underline{\text{var}} := E$

$E \cdot \text{exptype} \leftarrow \text{var.type}$ .

$E_0 \rightarrow E_1 + E_2$

$E_0 \cdot \text{type} \begin{cases} \text{int} & \text{if } E_1 \cdot \text{type} = \text{INT} \text{ AND} \\ & E_2 \cdot \text{type} = \text{INT} \\ \text{float} & \text{otherwise.} \end{cases}$

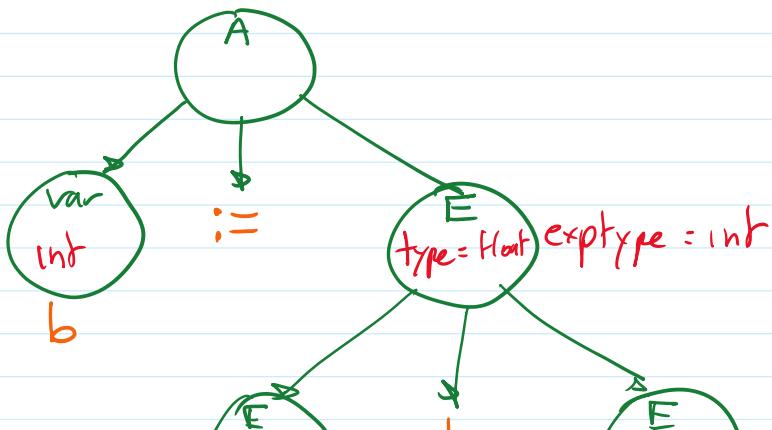
$E \rightarrow \underline{\text{var}}$

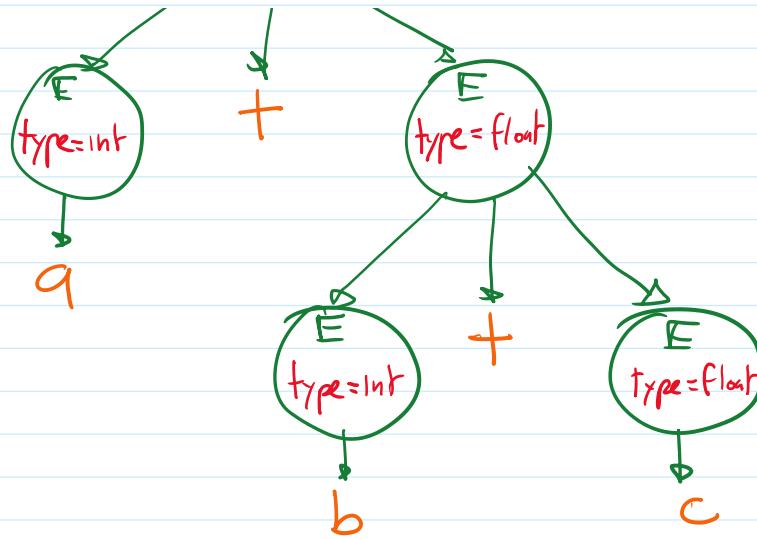
$E \cdot \text{type} \leftarrow \text{var.type}.$

$b := a + b + c$

Symbol table

a	int
b	int
c	float





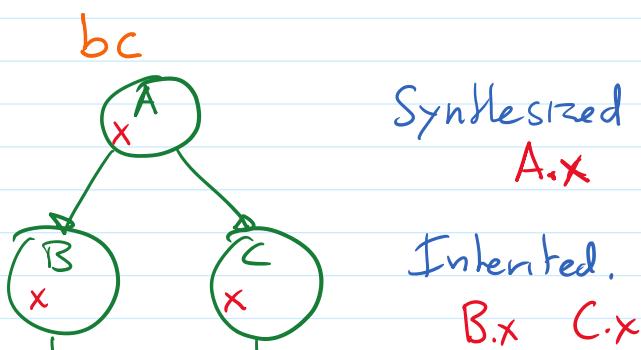
### COMPUTING ATTRIBUTE VALUES:

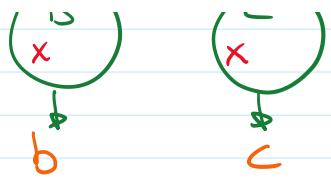
- Synthesized Attributes.
  - traverse the tree bottom up.
- Inherited Attributes
  - traverse the tree top to bottom
- Both kinds:
  - Multiple traversals bottom-up & top-down

### E.6 : Degenerate attribute grammar:

Attributes  $x$

$$\begin{array}{l}
 A \rightarrow BC \\
 B \rightarrow b \\
 C \rightarrow C
 \end{array}
 \quad
 \begin{array}{l}
 A.x \leftarrow C.x \\
 B.x \leftarrow A.x \\
 C.x \leftarrow B.x
 \end{array}$$





in revised.  
B.x C.x

