

SEMANTICS

The study of meaning.

e.g C++

```
float x, y, z;
```

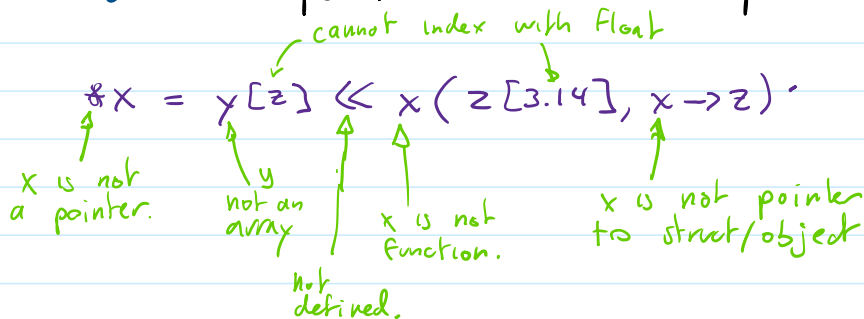
```
*x = y[z] << x(z[3.14], x->z);
```

follow the syntax rules.

but

it is **meaningless**.

because operators don't match operands.



In programming languages, semantics is concerned about identifiers and their correct use.

• TYPE CHECKING:

↳ Are named entities used in a manner consistent with their definition?

- using a variable with appropriate operators.
- ie: • call a function with appropriate number and types of parameters.
- members of compound types are used correctly

↳ How?

- The compiler/parser needs to remember the declaration of every named entity

• THE SYMBOL TABLE:

A table of named entities and their attributes.

names	attributes
x	int

examples of symbol table entries

• variable

float y; y : float

• constant

const int z = 3; z : int, 3, const.

• type

struct cell {
 int row, col;
}; cell : struct, 2, int row, int col

• function

void foo (int r, int c)
{
 ≡
}
foo : function, void, 2, int r, int c,

Entries are different in different programming languages.

C++:

```
int z[3];
float foo ( int x, bool y, char& c);
```

name	attributes
z	int array
foo	function 3 (int, bool, char&) <u>"float."</u>

Python

```
z = [1,2,3]
def foo ( x, y, c ) :
    .....
```

name	attr
z	[1,2,3]
foo	function 3

Pascal

```
z : ARRAY [10..12] OF INTEGER;
FUNCTION foo ( x : INTEGER,
               y : BOOL,
               VAR c : char ) : REAL
```

name	attributes
z	ARRAY, INT, 10..12
foo	function 3 Integer : REAL. Bool VAR char.

C++ Consequences.

```
int foo ( string y );
```

foo (string)

```
int foo ( char* y );
```

foo(char*)

```
string& foo ( string y );
```



string& foo (string y);

Pascals enumerated Arrays:

```
days = {Monday, Tuesday, Wednesday, Thursday, Friday}
z : ARRAY [Monday..Friday] OF INTEGER;
```

• TYPE CONVERSION:

Some languages will automatically apply type conversion:

e.g. $a+x$ $a: float$
 $x: int$

most parsers will identify the need for conversion and automatically convert.

$int \rightarrow float$
 $float \rightarrow int$

What conversions are Implicit?

C++ flexible

Python runs and maybe crash.

Pascal Strict.

E.g. C++ flexibility

$int \leftrightarrow float \leftrightarrow char \leftrightarrow bool$
↑ ↑
all sizes all-sizes

```
foo ( Cat c )
{
}
}
```

foo(3);

```
class Cat
{
    void meow();
    Cat ( int h )
    {
    }
}
```

becomes an implicit conversion unless specified explicit.

E.g Pascal is strict.

C

Pascal.

↳ Pascal is strict.

C

```
int x = 70;
char c = '!';
```

```
while ( ('W' - x) ) {
  c = '!' + x;
  print("%c", c);
  x = x + 1;
}
```

0 is false.
everything else
is true

Pascal.

```
VAR x : INTEGER = 70;
    c : CHAR = '!';
```

```
WHILE ORD('W') - x > 0 DO BEGIN
```

```
  c := CHR( ORD('!') + x );
```

```
END;
```

explicit conversions.

• TYPE EQUIVALENCE:

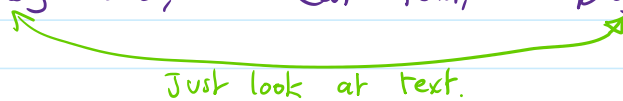
$$a = b$$

↳ when are two entities of the same type?

1) Name type equivalence:

- two entities are of the same type if they are declared using the same type name.

```
Dog bob;    Cat tom;    Dog fido;
```



1) Structural type equivalence:

- two entities are of the same type if they share the same internal structure.

Imagine:

C²

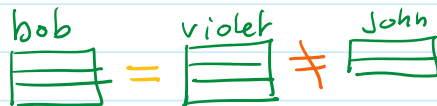
```
struct pnt
{
  int x,y,z;
}
```

```
struct color
{
  int r,g,b;
}
```

```
struct cell
{
  int r,c;
}
```

Under structural type equivalence, which are of the same type?

```
pnt bob
color violet
cell john
```



bob = violet;

• Hard to implement:

- you need to test structure
- ok... 1... 1... 1... 1... 1... 1... 1... 1... 1... 1...

- Hard to implement:
 - you need to test structure
 - structure can be nested.

• STRONG vs WEAK TYPING:

A characterization of Programming Languages

"Strongly" typed if

- type violations are detected at compile/parse time.
- type conversions are explicit
- the type of a named variable remains fixed.

General idea:

Detect errors at compile time,
instead of letting them happen at runtime.

