

15 Static and Dynamic Typing and Scope

Friday, November 1, 2024 12:42 PM

"Static" :- as in the text.

"Dynamic" :- as the program runs.

- TYPE BINDING:

¿ When does a variable gets its type?

- Static type Binding

Type stated somewhere in the text

- Explicit Statement

C++ `int x; string s;`

Pascal `VAR
 X: INTEGER;
 S: STRING;`

- Implicit Declaration

FORTRAN `I J K` are integers
 everything else is float
 (unless otherwise specified)

PERL
 name carries type.

`@x` array

`%x` hashtable / Dictionary

`$x` scalar

- Static Type Prediction

C++ `auto`

type is deduced from initialization.

Go `VAR float x` ...

Go
VAR float x
:
x = 3.14

VAR x := 3.14
↳ use type deduction

- Dynamic type Binding

- type is bound when a value is assigned as the program runs

- Type can change, by another assignment.

Python, JavaScript, Ruby, Lua

- SCOPE:

- The scope of a variable, is the range of statements over which the variable is visible.

- Global

BASIC

- Static Scope:

- range is determined by program text

C / C++

- scope is based on "blocks"

{ block }

e.g

int z; ← Global

int foo ()

{
int z; ← local variable masks global variable.

{
int z; ← inner scope masks outer scope.

z = 3 + 7;

}
}
}

Pascal scope is Modul or Function based.

```
VAR
  z: INTEGER;
```

```
FUNCTION foo(): INTEGER
```

```
VAR
  z: INTEGER
```

```
BEGIN
```

```
  z := 3+7;
```

```
END.
```

} the scope of z is the whole function.

• How does the Symbol table handles nested scopes

- Spit symbol table into "frames"
- As program is parsed:
 - new scope.- push a new frame
 - end scope.- pop a frame
- Symbol table is a stack.

e.g. C++

```
int x;
void foo ( int y )
{
  x = y * 10;
}

int main()
{
  int x, y;
  cin >> x >> y;
  if ( x < y ) {
    int x, z;
    cin >> z;
    x = y + z;
  }
  for( int x=0; x<y; x++) {
    cout << x;
  }
  cout << x;
}
```

Symbol Table	
x	int
foo	(int)
main	()

y	int

x	int
y	int

x	int
z	int

x	int

- Dynamic Scope.

- used in some esoteric programming Languages

Lisp, Bash

- '70 important among the Lisp community

- When is a variable visible? Depends on program execution!

- if fun1() calls fun2() then the variables of fun1() are available in fun2()

E.G. Imagine C++ w/ Dynamic Scope.

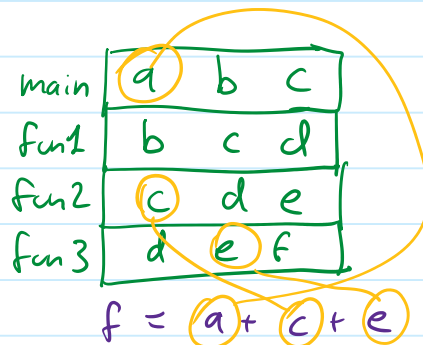
```
void main ()
{
  int a, b, c
  ... ..
}
```

```
void fun1 ()
{
  int b, c, d
  ... ..
  b = a + 1
}
```

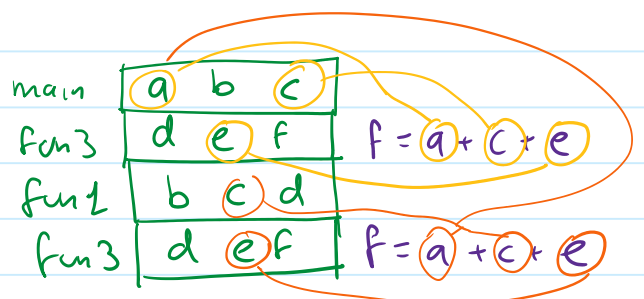
```
void fun2 ()
{
  int c, d, e
  ... ..
}
```

```
void fun3 ()
{
  int d, e, f
  ... ..
  f = a + c + e
}
```

Suppose
main()
└ fun1
└└ fun2
└└└ fun3



Suppose
main
fun3
└ fun1
└└ fun3



—●—●— EOF