

12 Indirect Left Recursion Elimination

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Eg 0

$A \rightarrow Ba \mid d$ $\{d, dcba, dcba cba, dcba cba cba, \dots\}$

$B \rightarrow Cb$

$C \rightarrow Ac$

parse-A()
 ↳ parse-B()
 ↳ parse-C()
 ↳ parse-A()
 ↳

A
Ba
Cba
Acba
dcba

• Remedy: Indirect Left Recursion Elimination

The Algorithm:

notation

A_i : non-terminal

$\delta \gamma$: Sequences of terminals & non-terminals

• Enumerate all non-terminals A_1 to A_n

• FOR $i \leftarrow 1$ TO n DO

 FOR $j \leftarrow 1$ TO $i-1$ DO

 Let A_j productions be

$A_j \rightarrow \delta_1 \mid \delta_2 \mid \delta_3 \mid \dots \mid \delta_k$

 Replace each production of the form

$A_i \rightarrow A_j \gamma$

 by $A_i \rightarrow \delta_1 \gamma \mid \delta_2 \gamma \mid \dots \mid \delta_k \gamma$

 Eliminate Direct left recursion in A_i

 // Note: This will insert a new rule in our sequence.

$A_1, A_2, A_3, A_4, \dots, A_n$

i

j

Example of inner loop:

$A_j \xrightarrow{\delta_1} a \mid \xrightarrow{\delta_2} b$

$A_i \rightarrow A_j x \gamma$

$A_i \rightarrow ax \mid bx$

Trace #1

$i=1$ A

1 $A \rightarrow Ba \mid d$
2 $B \rightarrow Cb$
3 $C \rightarrow Ac$

A
Ba
Cba

i=1 A
edlr σ

i=2 B
j=1 A
edlr σ

i=3 C
j=1 A
C $\rightarrow$ Bac | dc
j=2 B
C $\rightarrow$ Cbac | dc
edlr
C $\rightarrow$ dcC'
C' $\rightarrow$ bacC' | λ

i=4 C'
j=1 = A
j=2 = B
j=3 = C
edlr

1 A $\rightarrow$ Ba | a
2 B $\rightarrow$ Cb
3 C $\rightarrow$ Ac



A $\rightarrow$ B a | d
B $\rightarrow$ C b
C $\rightarrow$ dcC'
C' $\rightarrow$ bacC' | λ

A
Ba
Cba
dcC'ba
dcba
dcba
dcba

Ba
Cba
Acba
dcba
A
Ba
Cba
dcC'ba
dcba

Trace #2

1 S $\rightarrow$ Sx | Ay | Az
2 A $\rightarrow$ Sa | Bb | λ
3 B $\rightarrow$ Bp | Sq | r



A_i = S
edlr
S $\rightarrow$ AyS' | AzS'
S' $\rightarrow$ xS' | λ

A_i = S'
A_j = S
edlr

A_i = A
A_j = S
A $\rightarrow$ AyS'a | AzS'a | Bb | λ
A_j = S'
edlr
A $\rightarrow$ BbA' | A'
A' $\rightarrow$ yS'aA' | zS'aA' | λ

• FOR i ← 1 to n DO

FOR j ← 1 to i-1 DO

Let A_j productions be

A_j $\rightarrow$ δ_1 | δ_2 | δ_3 | ... | δ_k

Replace each production of the form

A_i $\rightarrow$ A_j γ

by A_i $\rightarrow$ $\delta_1 \gamma$ | $\delta_2 \gamma$ | $\delta_3 \gamma$ | ... | $\delta_k \gamma$

$$A \rightarrow BbA \mid A'$$

$$A' \rightarrow yS'aA' \mid zS'aA' \mid \Lambda$$

$$A_i = A'$$

$$A_j = S$$

$$A_j = S'$$

$$A_j = A$$

$$\text{edlr}$$

$$A_i = B$$

$$A_j = S \quad B \rightarrow Bp \mid \underline{A}yS'q \mid \underline{A}zS'q \mid r$$

$$A_j = S'$$

$$A_j = A \quad B \rightarrow Bp \mid \underline{BbA'}yS'q \mid \underline{A'}yS'q \mid BbA'zS'q \mid \underline{A'}zS'q \mid r$$

$$A_j = A'$$

$$B \rightarrow \underline{Bp} \mid \underline{BbA'}yS'q \mid yS'aA'yS'q \mid zS'aA'zS'q \mid yS'q \mid \underline{BbA'}zS'q$$

$$\mid yS'aA'zS'q \mid zS'aA'zS'q \mid zS'q \mid r$$

$$\text{edlr} \quad B \rightarrow yS'aA'yS'qB' \mid zS'aA'yS'qB' \mid yS'qB' \mid yS'aA'zS'qB' \mid zS'aA'zS'qB' \mid zS'qB' \mid rB'$$

$$B' \rightarrow pB' \mid bA'yS'qB' \mid bA'zS'qB' \mid \Lambda$$

$$A_i = B'$$

$$A_j = S$$

$$A_j = S'$$

$$A_j = A$$

$$A_j = A'$$

$$A_j = B$$

$$\text{edlr} \quad \emptyset$$

$$S \rightarrow Sx \mid Ay \mid Az$$

$$A \rightarrow \underline{S}a \mid Bb \mid \Lambda$$

$$B \rightarrow Bp \mid S_q \mid r$$

$$S \rightarrow AyS' \mid AzS'$$

$$S' \rightarrow xS' \mid \Lambda$$

$$A \rightarrow BbA' \mid A'$$

$$A' \rightarrow yS'aA' \mid zS'aA' \mid \Lambda$$

$$B \rightarrow yS'aA'yS'qB' \mid zS'aA'yS'qB' \mid yS'qB' \mid yS'aA'zS'qB' \mid zS'aA'zS'qB' \mid zS'qB' \mid rB'$$

$$B' \rightarrow pB' \mid bA'yS'qB' \mid bA'zS'qB' \mid \Lambda$$

—o—o

Replace each production of the form
 $A_i \rightarrow A_j \gamma$
 by $A_i \rightarrow \delta_1 \gamma \mid \delta_2 \gamma \mid \dots \delta_k \gamma$
 Eliminate direct left recursion in A_i