

Language Exploration : Lisp

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Language Exploration

object oriented

Simula
Smalltalk
Java

C++ Python

C Fortran
Pascal

Imperative

Functional

LISP

Clojure

Haskell

OCaml

• FUNCTIONAL PROGRAMMING

1.- No assignment. (Just labels)

2.- Computation is performed by
"function composition"

- No sequential execution
- No iteration, Recursion

$h(f(g(x)))$

3.- Functions are "first-class"

- can be passed as arguments to functions.
- can be return values from functions.

• LISP

List Processor.

1960's by John McCarthy

- While working in the "lambda calculus" of Alonzo Church.
concept:- Anonymous functions.

$f(x)$

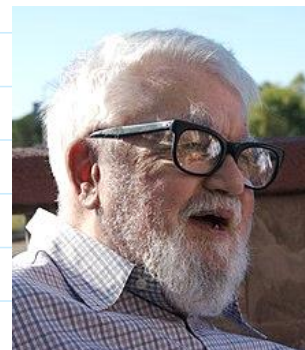
$\sin(y)$

$f.x$

$\sin.y$

$\lambda.x$

$\lambda.y$



- The CORE of LISP

Syntax:

- atom:- a sequence of letters, numbers or symbols (ex: ; # ' \)

e.g. foo apple 123 cat2
the-snake is-even? a-b

- expressions:- -an atom

or

-a list of zero or more expressions separated by spaces enclosed in parentheses.

e.g. foo (cat2 123) ()
(a b c) (foo (bar) (apple 123))
(apple)

Semantics.

(a b c d)

data

function application
apply function a to parameters b c d

e.g. (sin 4.57)

(fib 7)

(gcd 37 27)

- Basic Forms (built in functions)

- (quote x) evaluates to x

e.g.

(quote a) $\Rightarrow$ a

'a $\Rightarrow$ a

(quote (a b c)) $\Rightarrow$ (a b c)
'(a b c)

- Special atoms.

t - true

nil } false
()

- $(\text{atom } x)$:- t if x evaluates to an atom
nil otherwise.

eg. $(\text{atom } 'apple) \Rightarrow t$
 $(\text{atom } '(a\ b\ c)) \Rightarrow \text{nil}$
 $(\text{atom } (a\ b\ c)) \otimes$ unknown function a
 $(\text{atom } (\text{atom } 'a)) \Rightarrow (\text{atom } 't) \Rightarrow t$
 $(\text{atom } (\text{atom } 'a)) \Rightarrow \text{nil}$
 $(\text{atom } 'nil) \Rightarrow t$

- $(\text{eq } x\ y)$:- t if x and y evaluates to the same atom or both the empty list.

eg. $(\text{eq } 'a\ 'a) \Rightarrow t$
 $(\text{eq } 'a\ 'b) \Rightarrow \text{nil}$ $(= 'a\ 'a)$

- $(\text{car } x)$:- expects x to be a list
evaluates to the first element of x

eg. $(\text{car } '(a\ b\ c)) \Rightarrow a$
 $(\text{atom } (\text{car } ((\text{eq } 'a\ 'b)\ \text{nil}))) \Rightarrow$
 $(\text{atom } (\text{car } (\text{nil}\ \text{nil}))) \Rightarrow$
 $(\text{atom } \text{nil}) \Rightarrow$
 t

- $(\text{cdr } x)$:- expects x to be a list
evaluates to the list of elements after the first element of x

eg. $(\text{cdr } '(a\ b\ c)) \Rightarrow (b\ c)$
 $(\text{cdr } '(a)) \Rightarrow ()$
 $(\text{cdr } (\text{cdr } '(a\ b\ c))) \Rightarrow (c)$

- $(\text{cons } x\ y)$ expects y to be a list
evaluates to the list that consists of x followed by y

$(\text{cons } 'a \ '(b\ c)) \Rightarrow (a\ b\ c)$
 $(\text{cons } 'd\ \text{nil}) \Rightarrow (d)$

- Conditional Form:

$(\text{cond } (p_1\ e_1) (p_2\ e_2) (p_3\ e_3) \dots (p_n\ e_n))$

the p expressions are evaluated in order until one is t , then, the corresponding e expression is evaluated, the result is the value of the cond expression

if no p expression evaluates to t , cond returns nil

e.g.

$(\text{cond } ((\text{eq } 'a\ 'b) \text{'first})$
 $((\text{atom } 'a) \text{'second})) \Rightarrow \text{'second.}$

shorthand $(\text{if } p\ e_1\ e_2) \quad (\text{cond } (p\ e_1)$
 $(t\ e_2))$

- Functional Forms.

- $(\text{lambda } (a_1\ a_2\ a_3 \dots a_n) e)$

an anonymous function with parameters $a_1, a_2 \dots a_n$ and expression e as its body

$(\text{lambda } (a\ b\ c) (\text{cons } a (\text{cons } b (\text{cons } c\ \text{nil}))))$

- $(\text{funcall } f\ (a_1\ a_2 \dots a_n))$ expects f to be a function applies f to parameters $a_1 \dots a_n$

e.g.

$(\text{funcall } (\text{lambda } (a\ b\ c)$
 $(\text{cons } a (\text{cons } b (\text{cons } c\ \text{nil}))))$
 $(\text{'apple } \text{'banana } \text{'orange}))$
 $\Rightarrow (\text{apple } \text{banana } \text{orange})$

- $(\text{defun } \text{foo } (a_1\ a_2 \dots a_n) e)$

defines a named function with name foo and parameters $a_1 \dots a_n$ and body e .

and parameters $a_1 \dots a_n$ and body e .

e.g.

```
(defun make-triple (a b c)
  (cons a (cons b (cons c nil))))
```

```
(make-triple 'apple 'banana 'orange)
⇒ (apple banana orange)
```

- `(eval e)` evaluates expression e .

- Special Atoms / Functions

numbers are understood as themselves.
+ - * / mod div.

```
(+ 5 3 7)
```

- Demo Functions

- `is-null?`

- `Fizz buzz`

3, not 5	Fizz
5, not 3	Buzz
both	Fizzbuzz.

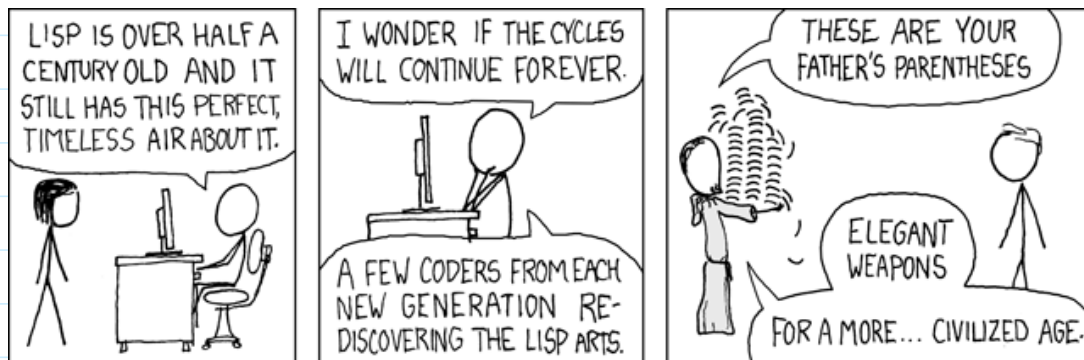
```
(+ 1 5 7)
```

```
(defun mysum. (L) ~~~~~)
```

```
(defun mycount (x L) ~~~~~)
```

```
  L
  () ⇒ nil
  (x ~~~~~) ⇒ 1 + (mycount x cdr(L))
```

$l() \Rightarrow nil$
 $(x \text{ ~~~~~}) \Rightarrow 1 + (mycount \times cdr(L))$
 $(y \text{ ~~~~~}) \Rightarrow (mycount \times cdl(L))$



$n \text{ is even} \begin{cases} \text{true} & \lambda.n \quad n/2 \\ \text{False} & \lambda.n \quad n*3 + 1 \end{cases}$

31 99 47 142 71 1

$1 = (1)$